

(12)

EUROPEAN PATENT APPLICATION

(43) Date of publication:
03.04.1996 Bulletin 1996/14

(51) Int Cl.⁶: H04L 29/00

(21) Application number: 95306439.1

(22) Date of filing: 14.09.1995

(84) Designated Contracting States:
DE FR GB NL

(30) Priority: 20.09.1994 US 309213

(71) Applicant: AT&T Corp.
New York, NY 10013-2412 (US)

(72) Inventors:

- Netravali, Arun Narayan
Westfield, New Jersey 07090 (US)
- Sabnani, Krishan Kumar
Westfield, New Jersey 07090 (US)

(74) Representative:

Watts, Christopher Malcolm Kelway, Dr. et al
AT&T (UK) Ltd.
5, Mornington Road
Woodford Green Essex, IG8 0TU (GB)

(54) Computer network gateway

(57) A gateway system for connecting first and second communication networks has first and second network end nodes and a protocol converter. The gateway system operates to provide input-output services of both network protocols to users of both networks in communicating through the gateway system. The gateway system employs associated complement protocol circuits in the end nodes and the protocol converter to provide these services. An efficient technique for generating the complement protocol of a first network is to prune finite state machines representing the operation of the second network protocol to provide the services of the second protocol that are not provided by the first protocol. The complement protocol circuits for the second network are then constructed based on the pruned finite state machines representing the operation of the first protocol.

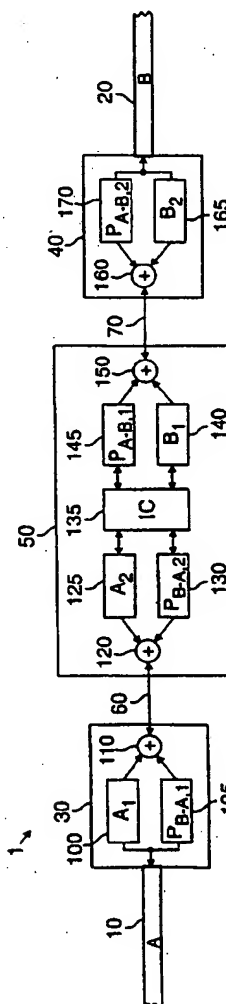


FIG. 2

Description

Field of the Invention

This invention relates to gateways between computer networks, and more particularly to the synthesis of a gateway system that permits users of the associated networks to transmit and receive information through the gateway system using services available on either network.

Background of the Invention

Computer networks and other data communication networks that employ communication protocols, such as Systems Network Architecture (SNA), DECNET, and Open Systems Architecture (OSI), are currently utilized around the world. A network protocol is a well defined set of rules governing the syntax, semantics and timing of the data transmission on the network. With the widespread use of electronic mail and file transfer services, it has become increasingly important to provide interconnections between networks. Three commonly used types of network interconnections are bridges, routers and gateways.

Conventional network interconnection systems typically have been built using *ad hoc* techniques. Problems in automatically synthesizing interconnections include overcoming network architectural mismatches and designing the protocol converters between the networks. Architectural mismatches involve differences in network layer functionalities, bit rates, buffering, interfaces and hand-off procedures. Since considerable judgment, trade-off, analysis, and decision making is required to overcome architectural mismatches, automatic generation of systems to overcome the architectural mismatches is not foreseeable in the near future.

Protocol converters translate messages from one protocol to another to rectify any mismatches between network protocols. Recently, techniques have been developed that automatically generate a protocol converter from formal specifications of the corresponding network protocols. Known automatic generation techniques include those described in S. Lam, "Protocol Conversion," *IEEE Trans. Soft. Engr.*, Vol. 14, No. 3, pp. 353-362, March 1988 (an earlier version was presented at SIGCOMM '86); and Y.W. Yao, W.S. Chen, and M.T. Liu, "A Modular approach to constructing Protocol Converters," *Proc. INFOCOM '90*, pp. 572-579, June 1990.

Automatic generation of protocol converters based on formal specifications has the following advantages: (1) the operation of the protocol conversion can be proved correct; (2) the protocol converter can be constructed in a relatively short period of time; and (3) the protocol converter can be changed quickly to adapt to changes in the protocol specification. However, most currently existing automatic synthesis techniques are relatively complex and inefficient due to a lack of effective

procedures that can handle large realistic protocols, as well as due to the difficulty in preparing the formal specifications of the protocols.

However, an example of an efficient automatic synthesis technique is described in D.M. Kristol, D. Lee, A.N. Netravali, and K. Sabnani, "A Polynomial time Algorithm for computing Protocol Converters from Formal Specifications," *IEEE/ACM Transactions on Networking*, pp. 217-229, April 1993, and in U.S. Patent No. 5,327,544, which is assigned to the assignee of the present invention. Both of these references are hereby incorporated by reference. The automatic synthesis technique described in these references generates a protocol converter that provides the greatest common subset of services available on the associated network protocols. This common subset of services is available to users that transmit and receive information through the protocol converter. Typical network services include data transfer, database inquiries and file transfers.

However, this technique is not transparent to users on either network because all the services available to the users on that network may not be part of the implemented subset in the protocol converter, and thereby not useable when communicating through the protocol converter. Therefore, a need exists for a method of automatically generating a gateway system having a protocol converter that provides a super-set of services to network users that transmit and receive information through the gateway system.

Summary of the Invention

A gateway system for connecting first and second communication networks according to one embodiment of the invention may include first and second network-end nodes and a protocol converter. Such a gateway system employs associated complement protocol circuits in the end nodes and the protocol converter to provide those services of one protocol that are not provided by the protocol of the corresponding network attached to that end node. As a consequence, the gateway system provides a super-set of services of both network protocols to users of both networks in communicating through the gateway system.

A method for generating a gateway system between first and second networks according to an embodiment of the invention first determines a common subset of services equal to an intersection of services provided by the protocols. Finite state machines ("FSM") are then generated which characterize first and second complement protocols of those services provided by the network protocols. The first complement service FSM contains those services provided by the second network protocol that are not provided by the first network protocol. Likewise, the second complement service protocol contains those services provided by the first network protocol and not the second network protocol.

Sets of intercommunicating finite state machines

representing each network protocol are then pruned to provide first and second complement network protocol finite state machines. The end nodes and protocol converter are then constructed using circuits that perform the operations of the sets of protocol finite state machines for the protocols, and the first and second complement network protocol finite state machines, as well as an interface converter. The method for generating the gateway system is simple, efficient, and can be accomplished in a polynomial-based number of steps.

The above-discussed features, as well as additional features and advantages of the present invention, will become more readily apparent by reference to the following detailed description and the accompanying drawings.

Brief Description of the Drawings

FIG. 1 is a schematic diagram of two data communication networks interconnected by a gateway system according to the present invention;

FIG. 2 is a detailed schematic diagram of a protocol converter and associated network end nodes in the gateway system of FIG. 1;

FIG. 3 is a flow diagram of the method for generating the gateway system of FIGS. 1 and 2;

FIG. 4 is a flow diagram of a method for pruning a protocol to generate a complement protocol for use in the method of FIG. 3;

FIG. 5 is a block diagram of a model arrangement of component finite state machines ("FSMs") to characterize a communication protocol that may be used in generating the gateway system of FIG. 1;

FIG. 6 is a transmitter FSM for use in a half-duplex alternating bit protocol (HABP) which may be implemented in one of the networks of FIG. 1;

FIG. 7 is a receiver FSM for use in the HABP;

FIG. 8 is a FSM for a first communications channel in the HABP;

FIG. 9 is a FSM for a channel communications channel in the HABP;

FIG. 10 is a FSM for a timer associated with the transmitter FSM of FIG. 6;

FIG. 11 is a service FSM for the HABP;

FIG. 12 is an arrangement of the FSMs for the HABP of FIGS. 6-11 according to the organization of FIG. 5;

FIG. 13 is a block diagram of an arrangement of component FSMs characterizing a full-duplex alternating bit protocol (FABP) which may be implemented in one of the networks of FIG. 1;

FIG. 14 is a transmitter FSM for use in the FABP;

FIG. 15 is a data retransmission timer FSM associated with the FSM of FIG. 14;

FIG. 16 is a receiver FSM for use in the FABP;

FIG. 17 is a sequence number storage FSM associated with the receiver FSM of FIG. 16;

FIG. 18 is a sequence number storage FSM associated with the transmitter FSM of FIG. 14;

FIG. 19 is a FSM for generating an acknowledgement signal associated with the receiver FSM of FIG. 16;

FIG. 20 is an acknowledgement timer associated with the FSM of FIG. 19;

FIG. 21 is a service FSM for the FABP;

FIG. 22 is a complement service generated according to the method of FIG. 3 based on the service FSMs of FIGS. 11 and 12;

FIG. 23 is a pruned transmitter FSM generated according to the method of FIG. 4 based on the FSM of FIG. 14;

FIG. 24 is a pruned data retransmission timer FSM generated according to the method of FIG. 4 based on the FSM of FIG. 15;

FIG. 25 is a pruned receiver FSM generated according to the method of FIG. 4 based on the FSM of FIG. 16;

FIG. 26 is a pruned sequence number storage FSM generated according to the method of FIG. 4 based on the FSM of FIG. 17;

FIG. 27 is a pruned sequence number storage FSM generated according to the method of FIG. 4 based on the FSM of FIG. 18;

FIG. 28 is a pruned FSM for generating an acknowledgement signal generated according to the method of FIG. 4 based on the FSM of FIG. 19; and

FIG. 29 is a block diagram of a gateway system between a HABP network and a FABP network employing the FSMs of FIGS. 14-28 arranged

according to the present invention.

Detailed Description

An embodiment of an interconnection of two networks by a gateway system 1 according to the present invention is shown in FIG. 1. In FIG. 1, first and second data communications networks 10 and 20 are connected to associated first and second network end nodes 30 and 40, respectively. The end nodes 30 and 40 are connected to a protocol converter 50 by first and second communications channels 60 and 70, respectively. The gateway system 1 consists of the end nodes 30 and 40, and the protocol converter 50.

Each one of the data communications networks 10 and 20 may be connected to respective computers 80, printers 90, file servers 100 and other data storage devices 110. Additional suitable data communication equipment may also be attached to the networks 10 and 20. In FIG. 1, the first data communications network 10 employs a protocol A and the second data communications network 20 employs a protocol B.

The operations of a network protocol may be modeled by a set of intercommunicating finite state machines (FSMs). Exemplary sets of FSMs characterizing half-duplex and full-duplex communication protocols are depicted in FIGS. 6-11 and 14-21, respectively, which are described in greater detail below. In the following description, interprocess input-output operation notation similar to that used in the language CSP is used to specify the intercommunication between protocol FSMs. CSP stands for Communicating Sequential Processes which is described in C.A.R. Hoare, "Communicating Sequential Processes," *Communications of the ACM*, Vol. 21, No. 8, pp. 666-677, August 1978 ("Communicating Sequential Processes").

According to this notation, a FSM sends a message to other FSMs by means of an "output operation" designated by the symbol "!". Where, for example, there are two FSMs, machine #1 and machine #2, an output operation in machine #1 which sends the message "msg" to machine #2 is denoted by *machine2!msg*. A FSM is also capable of receiving messages. For each message that is sent by one FSM there must be a receipt of the message by at least one other FSM. A FSM receives a message from another FSM by means of an "input operation" designated by the symbol "?". With respect to the output operation exemplified above, the corresponding input operation in machine #2 is denoted by *machine1?msg*, which means receive the message "msg" from machine #1.

The operations *machine1?msg* and *machine2!msg* are executed simultaneously because they represent different perspectives of the same event. Neither operation can be executed individually. Thus, if a first FSM attempts to do an output operation, it has to wait until a second FSM is ready to execute the corresponding input operation. The synchronized message exchange be-

tween two FSMs is called a *rendezvous*.

If the FSM name is not specified in an input or output operation, then the input-output operation can take place with any other FSM that is ready to execute the corresponding operation. For example, if a FSM has an operation *?msg*, then the receipt of "msg" from any other FSM will trigger the operation of the FSM. These types of input-output operations are often used to model a protocol's interactions with multiple users on multiple communication channels.

A typical FSM for use in an embodiment of this invention is a four-tuple $F = (\Sigma, V, p, s_0)$ where the symbol Σ is a set consisting of all the FSM's input-output operations and an internal operation; the symbol V is a finite set of states that the FSM may be in; the symbol p is a finite set of state transition functions, such as *machine!msg* or *machine?msg*; and the symbol s_0 is the initial state of the FSM F . While doing an internal transition from one state to another state, a FSM makes a state transition without interacting with any other FSM. For more information regarding FSMs see Z. Kohavi, *Switching Theory and Finite Automata Theory*, McGraw-Hill, pp. 275-315 (1978) and *Communicating Sequential Processes*.

A FSM may be represented as a directed graph (V, E) where the symbol V is the set of states in the FSM and the symbol E is the set of edges or possible state transitions between states. Each state of a FSM is represented in a directed graph by a circle encircling the designation of the state, such as the state 900 in the state machine *Htrans* of FIG. 6. Each edge is labeled by an input-output operation, belonging to the set Σ , which either triggers the state transition or is a result of the state transition. For example, edge 910 extending from state 900 in FIG. 6 is labeled with *?timeout*. Additionally, an edge may be labeled by two or more input-output operations which independently, or in conjunction, trigger the state transition, or are the result of it.

For notational purposes, an edge labeled by $a*b$ denotes an edge triggered by an input-output operation a followed by an input-output operation b . The symbol * represents the boolean AND operator. Therefore, a state transition, represented by an edge labeled by *?ack0*?cancel* is only triggered by the receipt of the messages "ack0" and "cancel". A state transition represented by an edge labeled by *?ack0!start* is triggered by the receipt of the message *ack0* and the sending of the message *start*. An edge labeled by $a + b$ denotes two edges that connect the same two states, one labeled by the input-output operation a and one labeled by the input-output operation b . The symbol $+$ represents the boolean OR operator. For example, a state transition represented by an edge labeled *?ack1 + ?start* is triggered by the receipt of either of the messages "ack1" or "start".

A FSM always starts in its initial state s_0 . The initial state is labeled 0 and may be additionally designated in the directed graph representing the FSM by the pres-

ence of a concentric circle surrounding the 0. When the FSM is in any given state it can execute any of the operations that label a transition from that state.

A protocol provides one or more services to the user of the protocol. These services are also represented as a FSM called the service FSM. However, there is not necessarily a one-to-one mapping or correlation between the set of services provided by a protocol and the set of FSMs that can represent the protocol operation. The input-output operations of the protocol represented by the service FSM are called service primitives of the protocol. According to international standards, a service primitive means an abstract message exchange between the protocol and its user or upper layer. An example of a service primitive is *!send(data)* where the primitive *send* is used to transmit a message *data*.

If an edge in one FSM is labeled by an input-output operation that has no corresponding operation in another FSM, the transition can never occur. For example, if an edge in machine #1 is labeled *machine2!msg*, and machine #2 contains no edge labeled *machine1?msg*, the transition in *machine1* can never occur. This characteristic of communicating FSMs is utilized in the method for generating a gateway system according to the present invention.

For any two FSMs F_1 and F_2 , a FSM designated $F_1 \# F_2$ can be built that corresponds to the joint behavior of F_1 and F_2 . The FSM $F_1 \# F_2$ is called a reachable FSM, or a composition of F_1 and F_2 . The process of constructing $F_1 \# F_2$ from the components F_1 and F_2 is called the reachability computation, or composing, and is well known in the art. Finding the reachable FSM is done by computing the reachable global states. A global state for $F_1 \# F_2$ is defined as a two-tuple (s_1, s_2) , where s_1 is the current state of F_1 and s_2 is the current state of F_2 . Computation of the reachable global states is described in G.v. Bochmann and C.A. Sunshine, "A Survey of Formal Methods," Computer Networks and Protocols, P.E. Green (Ed.), pp. 561-578, Plenum Press, May 1983.

A more detailed view of the gateway system 1 of FIG. 1 is shown in FIG. 2. In FIG. 2, the data communications network 10 is connected to a protocol circuit A_1 100 and a complement protocol circuit $P_{B-A,1}$ 105 in the network end node 30. Within the network end node 30, the circuits A_1 100 and $P_{B-A,1}$ 105 are connected to a first combined multiplexer-demultiplexer (MUX-DEMUX) 110.

The protocol circuit A_1 100 performs the operations specified in a portion of the set of protocol FSMs that characterize the protocol A and is described in greater detail below. The complement protocol circuit $P_{B-A,1}$ 105 performs operations that are specified in a portion of a set of complement protocol FSMs characterizing a complement protocol P_{B-A} . The complement protocol P_{B-A} provides those services of the protocol B implemented on the network 20 that are not provided by the protocol A of the network 10. The complement protocol circuit $P_{B-A,1}$ 105 and complement protocol P_{B-A} are also de-

scribed in greater detail below.

The MUX-DEMUX 110 is also connected to the first communications channel 60 which is connected to a second MUX-DEMUX 120 disposed in the protocol converter 50. Within the protocol converter 50, the MUX-DEMUX 120 is connected to an interface protocol circuit A_2 125 and an interface complement protocol circuit $P_{B-A,2}$ 130. The interface protocol circuit A_2 125 performs those operations characterized in the remainder portion, the set of FSMs for the protocol A that is not performed by the protocol circuit A_1 100, and is described in greater detail below. Likewise, the interface complement protocol circuit $P_{B-A,2}$ 130 performs those operations characterized in the remainder portion of the set of FSMs for the complement protocol P_{B-A} that are not performed by the complement protocol circuit $P_{B-A,1}$ 105. The interface complement protocol circuit $P_{B-A,2}$ 130 is also described in greater detail below. The MUX-DEMUXs 110 and 120 operate in conjunction with one another to multiplex and demultiplex signals on the communications channel 60 between the circuits A_1 100 and A_2 125, and between the circuits $P_{B-A,1}$ 105 and $P_{B-A,2}$ 130, respectively.

Within the gateway protocol converter 50, the circuits A_2 125 and $P_{B-A,2}$ 130 are connected to an interface converter IC 135. The interface converter IC 135 is further connected to the interface protocol circuit B_1 140 and an interface complement protocol circuit $P_{A-B,1}$ 145 which are described in greater detail below. The circuits B_1 140 and $P_{A-B,1}$ 145 are connected to a third MUX-DEMUX 150 which is connected to the second communications channel 70.

The communications channel 70 is further connected to a fourth MUX-DEMUX 160 contained in the network end node 40. Within the network end node 40, the MUX-DEMUX 160 is connected to the protocol circuit B_2 165 and a complement protocol circuit $P_{A-B,2}$ 170. The circuits B_2 and $P_{A-B,2}$ are also connected to the second data communications network 20.

The protocol circuit A_1 100 and the interface protocol circuit A_2 125 perform the communication functions of the protocol A as modeled by the corresponding set of FSMs. The determination as to which FSM of the set of protocol FSMs is to be implemented in which of the circuits A_1 100 or A_2 125 would be readily apparent to those skilled in the art. Generally, the functions specified by the FSMs of the set that are used to communicate with the users and upper-level of the network are implemented in the circuit A_1 100. Correspondingly, those functions characterized by the FSMs of the set that facilitate communication with other networks are implemented in the circuit A_2 125. An example of the arrangement of the functions of the various FSMs in the respective protocol circuits is described below with respect to FIG. 29.

The interface protocol circuit B_1 140 and the protocol circuit B_2 165 operate according to the protocol B as specified in the corresponding set of FSMs in an analogous manner to that of the circuits A_2 125 and A_1 100, respectively. Likewise, the functions of the complement

protocols P_{B-A} and P_{A-B} , which are characterized by the respective sets of FSMs are divided between the circuits $P_{B-A,1}$ 105 and $P_{B-A,2}$ 130, and between the $P_{A-B,1}$ 145 and $P_{A-B,2}$ 170, respectively. In a similar manner to that stated above, the complement protocol P_{A-B} characterizes the operations of those services provided by the protocol A that are not present in the protocol B.

By properly arranging the protocol and complement protocol functions in the end nodes 30 and 40, and the protocol converter 50, the gateway system 1 of FIG. 2 enables users attached to the network 10 to utilize all the services, or a super-set of the services, of the protocols A and B when communicating with the network 20. Likewise, users attached to the network 20 may utilize all the services of the protocols A and B when communicating with the network 10.

In operation, a service provided by both the protocols A and B is performed by a tandem combination of the exchange of corresponding service primitives between circuits A_2 125 and B_1 140 through the interface converter IC 135 in FIG. 2. In a similar manner, a service provided by the protocol A but not by the protocol B is performed by a tandem combination of the exchange of corresponding service primitives between circuits A_2 125 and $P_{A-B,1}$ 145 through the interface converter IC 135. A service provided by the protocol B but not by the protocol A is performed by a tandem combination of the exchange of corresponding service primitives between circuits B_1 140 and $P_{B-A,2}$ 130 through the interface converter IC 135.

Accordingly, the interface converter IC 135 identifies the service process that is attempting to communicate through the protocol converter 50. The interface converter IC 135 then routes the respective service primitives of that service between the appropriate circuits in the protocol converter 50. A suitable circuit for performing the operations of the interface converter IC 135 is a microprocessor or microcontroller connected to memory containing a mapping table for corresponding service primitives. An ASIC or PLA possessing the required routing capabilities may also be employed as the interface converter IC 135.

By utilizing the MUX-DEMUXs 110, 120, 150 and 160, the operations of the circuits A_1 100 and $P_{B-A,1}$ 105, A_2 125 and $P_{B-A,2}$ 130, B_1 140 and $P_{A-B,1}$ 145, and B_2 165 and $P_{A-B,2}$ 170, may be performed in parallel. One suitable technique for ensuring that messages will be exchanged only between protocol circuits, such as circuits A_1 100 and A_2 125, and between complement protocol circuits, such as circuits $P_{B-A,1}$ 105 and $P_{B-A,2}$ 130, is to use tags, such as a 1 or 2, to denote whether the message transmitted over the communications channel is from a respective protocol or complement protocol circuit. In an alternative embodiment, the MUX-DEMUXs 110, 120, 150 and 160 may be omitted if the corresponding protocol and complement protocol circuits are directly connected to one another.

The protocol circuit A_1 100 and the complement pro-

tol circuit $P_{B-A,1}$ 105, as well as the MUX-DEMUX 110 in the end node 30 may be arranged as separate or combined circuits. Further, the separate or combined circuits may be conventional microprocessors, ASICs or PLAs. The circuits of the protocol converter 50 and end node 40 may also be implemented as separate or various configurations of combined circuits in a similar manner.

FSMs characterizing the service operations provided to the users by the protocols A and B will be referred to by the symbols S_A and S_B , respectively. The input-output operations depicted by the FSMs S_A and S_B are the service primitives of the protocols A and B, respectively. The set of input-output operations of the circuits A_1 100 and A_2 125 with their respective local user or upper-levels will be referred to as IA_1 and IA_2 , respectively. An example of this type of input-output operation includes *deliver (data)* where a service primitive *delivery* is used to deliver a message *data*. In a similar manner, the set of input-output operations of the circuits B_1 140 and B_2 165 with their respective local users or upper-levels will be referred to as IB_1 and IB_2 , respectively.

For exemplary purposes, suppose the services provided by the protocols A and B are both data transfer protocols. In operation, the circuit A_1 100 may be given a service primitive for establishing a connection. In response, the circuit A_1 100 would generate a corresponding output operation to the circuit A_2 125. The circuit A_2 125 would generate a service primitive indicating to its local user, the interface converter IC 135, that a remote user to the circuit A_1 100 wishes to establish a connection. Upon receiving this service primitive, the interface converter IC 135 would generate an input service primitive for the protocol B indicating that a local user, the circuit A_2 125, wants to establish a connection.

Accordingly, the interface converter IC 135 would operate to cause the circuit B_1 140 to be the local user of the circuit A_2 125 by routing the output service primitives of the circuit A_2 125 to the input service primitives of the circuit B_1 140. The circuit B_1 140 would then generate a corresponding input operation to its local user, the circuit B_2 165 which provides communication with the network 20. Thus, a connection may be made by a user on the network 10 to the network 20 through the gateway system 1.

In the reverse direction, the output service primitives from the circuit B_1 140 are provided to the input service primitives of the circuit A_2 125 by the interface converter IC 135. However, since the networks protocols 10 and 20 are different, the input service primitives for the circuit A_2 125 may not have a one-to-one correspondence with the output service primitives of the circuit B_1 140. Therefore, a mapping or translation between the service primitives or input-output operations IA_2 and those of IB_1 must be implemented in the interface converter IC 135.

In most cases, this translation may be a direct mapping between two corresponding operation elements of the protocols. However, in other cases, the mapping may be more complex. For instance, if two output service

primitives x and y from the circuit A_2 125 are equivalent to a single input service primitive z for the circuit B_1 140, then the interface converter IC 135 must generate the output service primitive z upon receipt of either an input service primitive x or y . Particular service primitives of one protocol may perform functions that are not performed by any combination of service primitives of the corresponding protocol. The service primitives provided by these unmatched services are performed by the complement protocols P_{A-B} and P_{B-A} implemented in the circuits $P_{A-B,1}$ 145 and $P_{A-B,2}$ 170, and $P_{B-A,1}$ 105 and $P_{B-A,2}$ 130, respectively.

One suitable method 200 for generating the gateway system 1 and corresponding complement protocols P_{A-B} and P_{B-A} is illustrated in FIG. 3. Referring to FIG. 3, a FSM W representing the largest common subset of services of the protocols A and B is computed in step 210. One method for computing the FSM W is to compute the expression $W = S_A \# IC \# S_B$, where S_A and S_B are collections of FSMs that are pruned versions of the FSMs S_A and S_B , respectively. The pruned FSMs S_A and S_B are formed by removing the edges of the FSMs S_A and S_B for which there are no corresponding service primitives provided by the other service FSM S_B and S_A , respectively. The resulting FSM W may be reduced while maintaining its operational equivalence according to a method described in P.C. Kanellakis and S.A. Smolka, "CCS Expressions, Finite State Processes, and Three Problems of Equivalence," *Information and Computation*, Vol. 86, pp. 43-68 (1990), which is incorporated by reference herein.

Complement service FSMs S_{B-A} and S_{A-B} specifying service primitive operations for the complement protocols P_{B-A} and P_{A-B} are then generated in step 220. One method for generating the complement service FSMs S_{B-A} and S_{A-B} is to prune the service FSMs S_B and S_A with respect to the common services subset FSM W . According to this method, the complement service FSM S_{B-A} is generated by removing the edges of the service FSM S_B that correspond to present edges in the common subset FSM W .

After generation of the complement service FSMs S_{B-A} and S_{A-B} in step 220, the method 200 proceeds to step 230. In step 230, the complement protocol FSMs P_{B-A} and P_{A-B} are generated by pruning the set of protocol FSMs characterizing the protocols B and A , respectively, to provide only those services of the complement service FSMs S_{B-A} and S_{A-B} . One suitable technique for pruning the FSMs representing the protocols A and B is described in greater detail below with reference to FIG. 4.

The protocol converter 50 and network end nodes 30 and 40 are then constructed in step 240 according to the manner shown in FIG. 2. The protocol converter 50 and network end nodes 30 and 40 may each be comprised of at least one processor, such as a microprocessor, ASIC or PLA, to provide the operations of the corresponding interconnection of protocol and complement protocol FSMs in the circuits illustrated in FIG. 2. Meth-

ods for implementing FSMs in processors are well known in the art.

FIG. 4 depicts a suitable method 300 for pruning the protocol FSMs to generate the complement protocol FSMs P_{B-A} and P_{A-B} in step 230 of FIG. 3. The steps of the method 300 of FIG. 4 may be grouped as follows: those edges that are labeled with service primitives other than that found in the corresponding service FSM S are deleted in the sequence of steps 320-340; those edges of the protocol FSMs that have no matching service primitives after the first sequence of steps 320-340 are deleted in steps 350-400; and in steps 405-460, the method 300 computes and retains the strongly connected FSM components that start with the initial state, and discards the rest of the machine.

In FIG. 4, a data structure MATCH, a FSM list L and counter variables i, j, k, m are initialized to zero in step 305. Then, in step 310, the operations of a protocol P , such as protocol A used on the network 10 of FIG. 1, are characterized by a set of FSMs $F_i, i=1, \dots, N$, where the value N is the total number of FSMs in the set. The counter i is then incremented in step 320. In step 330, those state transitions in the FSM F_i that do not correspond to the service primitives utilized by the corresponding protocol service FSM S are deleted. If the protocol A of network 10 were used as the protocol P in the method 300, then the service FSM S_A will be used as the service FSM S in the method 300.

The counter i is then compared with the total number of FSMs N in step 340. If the counter i does not equal the total number of machines N , the method 300 repeats steps 320 and 330 to delete or prune state transitions having unmatched service primitives in the next FSM F_i of the set. However, if the counter i is equal to the total number of machines N , the method 300 proceeds to step 350 where the counter j is incremented.

After performing the step 350, the FSM F_j is evaluated in step 360 to determine if it contains edges corresponding to services primitives that are not matched by counterpart service primitives in any other FSM of the set of FSMs F . An edge representing a service primitive may be unmatched due to deletion of their corresponding edges in step 330. For example, if an edge of the FSM F_j has a label $!startx$ for sending a service primitive $startx$ and if the corresponding edge labeled $?startx$ in another FSM of the set was deleted in step 330, then the edge of the FSM F_j will be unmatched. If at least one edge of the FSM F_j is unmatched, the method 300 performs steps 370-390. In step 370, all identified unmatched edges of the FSM F_j are deleted, and in step 380, the FSM F_j is appended to the list L . In step 390, the counter m which corresponds to the length of list L is incremented. After performing step 390, the method 300 proceeds to step 400. However, if in step 360, it is determined that the FSM F_j contains only edges having matched service primitives, the method 300 will proceed directly to step 400.

In step 400, the counter j is tested and if it is not equal

to the total number of FSMs N in the set F , the method 300 repeats steps 350 and 360 to identify if the next FSM F_j contains edges having unmatched service primitives and amends the FSM F_j and list L accordingly in steps 370-390. Alternatively, if the method 300 determines, in step 400, that the counter j is equal to the total number of FSMs N then all the FSMs F of the set have been processed for edges having unmatched service primitives, and the method 300 proceeds to step 405.

In step 405, the method 300 determines whether the value of the counter m is zero. If the counter m has a value of zero, no FSM in the set of FSMs F had edges deleted in step 370. As a result, the method 300 proceeds directly to step 470 which is described below. However, if the counter m is not zero in step 405, the method 300 proceeds to step 410. In step 410, the counter k is incremented, and then, in step 420, a FSM F_k is removed from the list L for further processing by steps 430, 440 and 450. All FSMs contained in the corresponding list L have at least one edge deleted in step 370. The number of FSMs in the list L is represented by the counter m .

In step 430, a strongly connected component of FSM F_k that contains the initial state is determined and denoted C_k . In a strongly connected FSM component, each node has a directed path to every other node. In step 440, those edges of the FSM F_k that are not contained in C_k are deleted, and in step 450, the data structure MATCH is updated to reflect those edges deleted in step 440. The FSMs characterizing protocols that do not have infinite loops for data transfers, such as connection management and call setup protocols, may be made strongly connected by adding dummy transitions.

Then, in step 460, the method 300 determines whether the counters k and l are equal. If counters k and l are found equal, the method 300 proceeds to step 470. However, if the counters k and l are not found equal, the method 300 repeats steps 410-450 to increment the counter k and determine the strongly connected components for the remaining FSMs in the list L .

In step 470, the method 300 outputs the FSMs F_i , $i=1, \dots, N$, characterizing the pruned protocol P . The resulting FSMs F_i , $i=1, \dots, N$, are reduced component FSMs of the original set of FSMs that characterize the protocol P and which contain the services corresponding to those of the service FSM S . Accordingly, the complement protocol FSMs P_{B-A} and P_{A-B} can be generated by performing the method 300 twice, a first time with the protocol P being the protocol B and with the service FSM S being the service FSM S_{B-A} , and a second time with the protocol P being the protocol A and with the service FSM S being the service FSM S_{A-B} .

One feature of the method 300 is that it generates the pruned machines utilizing a polynomial, rather than an exponential, number of computational steps. An advantage of this feature is that it permits the pruning of relatively large and complex protocols, such as conventional computer network protocols, in a relatively short

computational time. Typical prior art pruning techniques require an exponential number of computational steps which is very impractical for the size and complexity of existing network protocols.

The data structure MATCH keeps track of the matching input-output services. Identical services may be associated with a counter value, which records their total number. Counter values of matching input-output services are associated with each other. When an input-output service is deleted, its associated counter value is decreased by one. When a counter value becomes zero, all the corresponding counter values are checked as follows: if a counter value μ has no matching counter value, i.e., all have become zero, then the counter value μ is changed to zero and all its associated edges in the FSM depicting that input-output service are deleted. The corresponding counter values of the counter value μ are processed similarly. This counter value updating is done iteratively until no changes have to be made to the involved counter values.

The total cost in the number of computations to ini-

tialize and to update the data structure MATCH is $O(\sum_{i=1}^N m_i)$, where m_i is the number of edges of F_i and where k is the number of FSMs in the set. A k -bit vector may be used to record whether the FSM F_i is in the list L . In one particular vector format, if the FSM F_i is in the list L , then the i -th bit of the vector is 1, otherwise it is 0. To update the vector and to check whether the FSM F_i is in the list L takes a constant time. Whenever a component FSM F_i is appended to the list L , at least one edge has been removed. Then, when the component FSM F_i is removed from the list L , the strongly connected component is constructed that contains the initial state. It takes time $O(m_i)$ to compute the strongly connected component of a FSM F_i , using a depth-first search. Since an FSM F_i can be removed from the list L only after it has been appended

to the list L , then the total computational cost is $O(\sum_{i=1}^N d_i m_i)$, where d_i is the number of edges deleted from the FSM F_i .

A synthesis of an exemplary gateway system 1 using the method 200 of FIG. 3 is described with respect to FIGS. 5-29. The two network protocols that are to be connected by the gateway system 1 in the example are a half-duplex alternating bit protocol ("HABP") and a full-duplex alternating bit protocol ("FABP"). The HABP transports data messages from a transmitter user process to a receiver process over a lossy communication channel. No data message is transported in the reverse direction for the HABP. Conversely, the FABP operates in a manner such that two remote processes are able to exchange data messages.

A general organization 500 of component FSMs

which characterize a protocol is illustrated in FIG. 5. In FIG. 5, the arrangement consists of end entities D₁ 510 and D₂ 520 which communicate with one another over communication channels C_{A12} 530 and C_{A21} 540. Each end entity D₁ 510 and D₂ 520 operates according to a respective group of the FSMs in the corresponding FSM set characterizing the protocol. The operation of the HABP may be characterized using five FSMs arranged according to the organization 500 of FIG. 5 as follows: a data transmitter machine *Htrans* shown in FIG. 6, which corresponds to the end entity D₁ 510 of FIG. 5, and which retransmits data received from a local user; a data receiver machine *Hrec*, shown in FIG. 7, which corresponds to the end entity D₂ 520, and receives the data retransmitted by the FSM *Htrans*; a forward communication channel machine C_{H12}, shown in FIG. 8, for transporting data messages from the FSMs *Htrans* to *Hrec* of FIGS. 6 and 7; a reverse communication channel machine C_{H21}, shown in FIG. 9, for transporting acknowledgement signals *ack0* or *ack1* confirming receipt of the data message; and a timer machine *Timer* shown in FIG. 10 which is a satellite FSM of the data transmitter machine *Htrans* and should also be contained in the end entity D₁ 510 of FIG. 5. Once started, the timer machine *Timer* of FIG. 10 will wait a predetermined time period and generate a time out signal *timeout* to the transmitter machine *Htrans* of FIG. 6, unless during that time period it receives a message *cancel* from the transmitter machine *Htrans*.

Further, a service FSM S_H for the HABP is shown in FIG. 11. In FIG. 11, the service FSM S_H consists of two states, 0 and 1. The FSM S_H depicts the data transfer operation of the HABP. An edge labeled *?dataI* denotes an input operation whereby the HABP receives a data message from the transmitter user process. A second edge labeled with *!dataO* represents an output operation whereby the HABP sends a data message to the receiver user process. FIG. 12 shows the five component HABP FSMs arranged according to the organization 500 of FIG. 5. Similar components in FIGS. 5 and 12 are like numbered for clarity, for example, the end entity 510.

In operation, the HABP, as represented by the FSMs *Htrans*, *Timer*, C_{H12}, C_{H21} and *Hrec* of FIGS. 6-10 and service FSM S_H of FIG. 11, starts with the transmitter function characterized by the FSM *Htrans* receiving a data message from a user using the input operation *?dataI*. The transmitter function FSM *Htrans* then sends the data message with a sequence number of either 0 or 1 to the receiver function, characterized by FSM *Hrec*, over the channel C_{H12}. The transmitter function FSM *Htrans* alternates the sequence numbers between 0 and 1 for each consecutive data message in order for the receiver function characterized by the FSM *Hrec* to identify whether a data message has been lost.

When the transmitter function FSM *Htrans* transmits the data message, it also starts a local timer operation characterized by the FSM *Timer*, and waits for the proper acknowledgement signal *ack0* or *ack1* over the channel

characterized by the FSM C_{H21}. Upon receiving the proper acknowledgement signal, the transmitter function cancels the timer by sending a message *cancel*. However, if no acknowledgement signal is received before the timer FSM *Timer* generates a signal *timeout*, the data message is retransmitted to the receiver FSM *Hrec*.

Upon receipt of a data message, the receiver function FSM *Hrec* determines whether it contains the proper sequence number and transmits the data message to the local user using the output operation *!dataO* of the service FSM S_H. The receiver function also sends the acknowledgement signal with a proper sequence number over the communications channel C_{H21} back to the transmitter function FSM *Htrans*.

FSMs characterizing the operation of the FABP may also be organized according to the system 500 of FIG. 5 as shown by the arrangement 700 in FIG. 13. In FIG. 13, the FABP is modeled using two identical protocol entities 710 communicating with one another by two communication channels C_{F12} and C_{F21}. For the FABP, each one of the protocol entities may consist of seven FSMs. The seven FSMs are: a transmitter function FSM *Ftrans*, shown in FIG. 14, for retransmitting outgoing data messages received from a user; a retransmission timer FSM *Retimer*, shown in FIG. 15, which operates in a similar manner to the FSM *Timer* of FIG. 10; a receiver function FSM *Frec*, shown in FIG. 16, which receives incoming data messages and acknowledgement signals; a sequence number storage FSM *NR*, shown in FIG. 17, for storing the next sequence number of expected message to be received by the receiver function FSM *Frec*; a sequence number storage FSM *NS*, shown in FIG. 18, for storing the sequence number of the next expected message to be sent by the transmitter function FSM *Ftrans*; an explicit acknowledgement FSM *Buf*, shown in FIG. 19, for sending acknowledgement signals confirming receipt of a data message; and an acknowledgment timer FSM *Atimer*, shown in FIG. 20, for performing the timing operations for the FSM *Buf*.

A service FSM S_F for the FABP is shown in FIG. 21. In FIG. 21, the service FSM S_F consists of four states 0, 1, 2 and 3. The FSM S_F depicts the data transfer operation of the FABP. During operation, the service FSM S_F receives a data message from one network user process by traversing edges that activate an input operation *?dataF*, and delivers it to the other user process by traversing the edges that activate an output operation *!dataF*.

Each end entity 710 of the FABP in FIG. 13 employs two methods of acknowledging receipt of data messages with the appropriate sequence number. In a first method, the receiver function FSM *Frec* uses the FSM *Buf* of FIG. 19 and associated FSM *Atimer* of FIG. 20 to generate acknowledgement signals *ack0* and *ack1*. In addition, a second method of acknowledging receipt of a data message is to piggy-back acknowledgement signals on return data messages.

A data message has a structure of *data[x][y]* where *data* is the information that is sent. The parameter *x* is

the sequence number of the message and may be either a 0 or 1. The parameter y is the piggyback acknowledgement, and may have a value of either a 0 or 1. The parameter y performs the same operation as the transmission of an acknowledgement signal $ackO$ or $ack1$ for a recently received data message. In operation, when a data message is received, the operation characterized by the FSM $Frec$ starts the operation represented by the FSM Buf which activates the process $Atimer$. If a data message is not sent out by the end entity 710 of FIG. 13 before the timer $Atimer$ expires, then the proper acknowledgement signal $ackO$ or $ack1$ is transmitted. The composition of these seven FABP protocol FSMs in FIGS. 14-20 represent a total of 3,192 states and 14,026 edges.

A gateway system generated according to the method 200 of FIG. 3 for connecting a HABP network and a FABP network would provide full duplex communication through the gateway between users on either network. For ease of illustration, the HABP will alternatively be referred to as the protocol A and the FABP as the protocol B in the following description. In generating such a gateway system according to step 210 of the method 200 of FIG. 3, a FSM W is computed corresponding to the largest common subset of services provided by the service FSM S_A , or the FSM S_H in this example, and the service FSM S_B , or the FSM S_F in this example. Since the service operations of the service FSM S_H of FIG. 11 are all included in the service FSM S_F , the resulting FSM W is equal to the service FSM S_H .

Then, according to step 220, complement service FSMs S_{A-B} and S_{B-A} are computed. The complement service FSMs S_{A-B} and S_{B-A} may be computed by removing those services provided by the FSM W from the service FSMs S_{A-B} and S_{B-A} , respectively. Accordingly, the complement service protocol FSM S_{A-B} is null because all the services provided by the FSM W are provided in the service protocol FSM S_H . However, the resulting complement service protocol S_{B-A} provides those services of the FSM S_F that transfer data in the reverse direction, or the direction not provided by the HABP. The resulting complement service FSM S_{B-A} is shown in FIG. 22.

The method 200 then proceeds to step 230 where the protocols A and B are pruned to form the complement protocol FSMs P_{A-B} and P_{B-A} . Since the complement service protocol FSM S_{A-B} is null the corresponding complement protocol FSM P_{A-B} is also null. However, pruning the protocol B, represented by the collection of FSMs of FIGS. 14-20, to provide those services of the complement service protocol S_{B-A} yields the FSMs illustrated in FIGS. 23-28. In FIGS. 23-28, the pruned FSMs are labeled by adding a prime to the name of the corresponding FSMs of FIGS. 14-20. For example, the transmitter FSM $Ftrans$ of FIG. 14 corresponds with the pruned transmitter FSM $Ftrans'$ of FIG. 23.

During the pruning step, the FSM $Atimer$ of FIG. 20 is pruned out of existence, while the sequence number

storage FSM NS of FIG. 18 and retransmission timer FSM $Retimer$ of FIG. 15 remain fully intact in forming the pruned counterparts in FIGS. 27 and 24. The pruned FSMs for the complement protocol P_{B-A} in FIGS. 23-28 provide only those data transfer functions in the reverse direction. The last step 240 of the method 200 is to construct the protocol converter 50 and network end nodes 30 and 40 according to the generated complement protocols in the manner shown in FIG. 2.

A gateway system 800 connecting a HABP network 10 and a FABP network 20 employing circuits according to the FSMs of FIGS. 23-28 is shown in FIG. 29. The circuits in FIG. 29 corresponding to those of FIG. 2 are like numbered for clarity, for example, the complement protocol circuit 105. In addition, those circuits in FIG. 29 which perform the operations of the FSMs illustrated in FIGS. 14-28 are like named for ease of illustration. For example, the circuit $Htrans$ in FIG. 29 and the corresponding transmitter function FSM $Htrans$ in FIG. 14.

In FIG. 29, an end node 30 of the network 10 contains the circuits A_1 100 and $P_{B-A,1}$ 105. The protocol circuit A_1 100 performs the data message transmission operation of the HABP of the network 10. The circuit A_1 100 contains the circuit $Htrans$ and the half duplex communication channels C_{H21} and C_{H12} which operate according to the FSMs of FIGS. 8 and 9, respectively. The complement protocol circuit $P_{B-A,1}$ 105 consists of circuits $Frec'$ and NS' for receiving data. The circuits A_1 100 and $P_{B-A,1}$ 105 transmit and receive messages from a protocol converter 50 through the MUX-DEMUXs 110 and 120 over the communication channel 60.

The protocol converter 50 of FIG. 29 includes the circuit A_2 125 for performing the operations of the HABP receiver FSM $Hrec$ of FIG. 7. The protocol converter 50 also consists of the interface complement protocol circuit $P_{B-A,2}$ 130, the interface converter IC 135, and the interface protocol circuit B_1 140 corresponding to the protocol end entity 710 in FIG. 13. The circuit B_1 140 performs the operations of the seven FABP protocol FSMs shown in FIGS. 14-20. The complement protocol circuits $P_{B-A,1}$ 105 and $P_{B-A,2}$ 130 in FIG. 29 operate in the manner specified by the complement protocol FSMs shown in FIGS. 23-28. The complement protocol circuits operations $Frec'$ and NS' for retransmitting data to a destination user and are, therefore, disposed in the circuit $P_{B-A,1}$ 105. Accordingly, the complement protocol circuit operations $Ftrans'$, Buf' , NR' and $Retimer'$ which receive the data from a local user have been disposed in the circuit $P_{B-A,2}$ 130.

The circuit B_1 140 of the protocol converter 50 is directly connected to a second network end node 40. The network end node 40 consists of the protocol circuit B_2 165 which performs the identical operation of the circuit B_1 140 in the protocol converter 50. Full-duplex communication between the circuits 140 and 165 occur over communication channels C_{F12} and C_{F21} . The circuit B_2 165 is also connected to the network 20.

The complement protocol circuits $P_{A-B,1}$ 145 and

P_{A-B,2} 170 and associated MUX-DEMUXs 150 and 160 of FIG. 2 are not used in the system 800 of FIG. 29 because no such circuits are required. Since the complement protocol FSM P_{A-B} and complement service FSM S_{A-B} are null, the corresponding complement protocol circuits are not required. Such circuits are not necessary in this example because there are no operations performed by the protocol HABP of network 10 that are not performed by the protocol FABP of the network 20. The circuits B₁ 140 and B₂ 165 perform identical operations in FIG. 29 because of the mirror-like processes of a full-duplex protocol. However, these circuits may perform different operations in other gateway systems 1 depending upon the protocols implemented on the networks 10 and 20.

In transmitting information from a network 10 user to a user process on the network 20, the information is transmitted through the circuit A₁ 100 in the network end node 30 to the circuit A₂ 125 in FIG. 29. The information from the circuit A₂ 125 is then directed to the circuit B₁ 140 by the interface converter IC 135. The information is then transmitted by the circuit B₁ 140 to the circuit B₂ 165 in the end node 40 where it is retransmitted to the user process on the network 20.

In the reverse direction, network 20 users may transmit data to network 10 users in FIG. 29, despite the fact that only an HABP is implemented in the network 10. In transmitting information from a network 20 user process to a network 10 user process, the information is transmitted through the circuit B₂ 165 in the network end node 40 to the circuit B₁ 140 in the protocol converter 50. The circuit B₁ 140 provides the information to the interface converter IC 135 which is configured to transmit the information to the circuit P_{B-A,2} 130. Then, the information is transmitted by the circuit P_{B-A,2} 130 to the circuit P_{B-A,1} 105 in the network end node 30. The circuit P_{B-A,1} 105 possesses the capability to retransmit the information to the user process on the network 10.

Therefore, the present invention provides a relatively short and simple technique for generating a gateway between different network protocols employing complement protocols which enables users of both networks to transmit and receive information between each other using a super-set of services provided by the protocols of both networks. One advantage of providing a super-set of services is that the gateway system between the networks may appear transparent to a user of one of the networks as he will be able to utilize all the services familiar to him on that network to exchange data with the other network.

In an alternative embodiment of the gateway system 1 of FIG. 2, complement protocols and corresponding circuits may be generated for only one of the networks using the method 200 of FIG. 3. For example, the complement protocol P_{B-A} and not P_{A-B} would be generated for such a system. Accordingly, the circuits P_{A-B,1} 145 and P_{A-B,2} 170 would be eliminated from the protocol converter 50 and end node 40. Likewise, the MUX-DEMUXs

150 and 160 may be removed as no messages need be multiplexed on the communications channel 70 between the circuits B₁ 140 and B₂ 165. As a result, users on the network 10 may utilize all the services of both of the protocols A and B when transmitting or receiving information using this alternative gateway system. However, users on the network 20 would be limited to using only those services common to protocols A and B when communicating through this alternative gateway system.

Although one embodiment of a method for generating a gateway system that provides a super-set of protocol services has been described in detail above, it would be readily understood by those having ordinary skill in the art that many modifications are possible in the described embodiment without departing from the present invention. All such modifications are intended to be encompassed by the claimed invention. For instance, a single microprocessor may be employed in the network end nodes and protocol converter to perform those operations specificized by the FSMs contained therein. Further, although the exemplary gateway system described above concerned computer networks, the present invention may be used for generating for providing gateway systems to other data and telecommunication networks.

Claims

1. A method for generating a gateway system between first and second networks implementing corresponding first and second communication protocols which provide first and second sets of services, respectively, wherein each protocol can be represented by a set of at least one finite state machine (FSM), the method comprising the steps of:

determining a common subset of services equal to a subset of the intersection of services provided by the protocols;

determining first and second complement service FSMs for the first and second network protocols, wherein the first complement service FSM contains those services provided by the second network protocol that are not characterized in the first network protocol, and wherein a second complement service FSM contains those services provided by the first network protocol that are not characterized in the second network protocol;

pruning the set of second network protocol FSMs to provide only operations corresponding to the services performed by the first complement service FSM to generate a first set of complement network protocol FSMs;

pruning the set of first network protocol FSMs to provide only operations corresponding to the services performed by the second complement service FSM to generate a second set of complement network protocol FSMs; and

constructing a gateway system based on the first and second set of protocol FSMs, the first and second set of complement protocol FSMs and an interface converter.

2. The method of claim 1, wherein each of the steps of pruning the set of protocol FSMs to generate a corresponding complement protocol FSM comprises the steps of:

deleting an edge from a first protocol service FSM corresponding to a service not characterized in a second protocol service FSM, and deleting an edge from the second protocol service FSM corresponding to a service not characterized in the first service FSM to form a first set of pruned FSMs;

deleting an edge from the first set of pruned FSMs which represent unmatched service primitives to form a second set of pruned FSMs; and

deleting an edge in the second set of pruned FSMs which are not within a strongly connected component of the corresponding second pruned FSM that contains the initial state to form a third set of pruned FSMs.

3. The method of claim 1, wherein the step of determining first and second complement protocols comprises the steps of:

pruning protocol FSMs characterizing the services provided by the protocols by removing edges in the service FSMs that correspond to services provided in the common subset of services.

4. A method for generating a gateway system between first and second networks implementing corresponding first and second communications protocols which provide first and second sets of services, respectively, wherein each protocol can be represented by a set of at least one finite state machine (FSM), the method comprising the steps of:

determining a common subset of services equal to a subset of the intersection of services provided by the protocols;

determining first and second complement service FSMs for the first and second network protocols, wherein the first complement service FSM contains those services provided by the second network protocol that are not characterized in the first network protocol, and wherein the second complement service FSM contains those services provided by the first network protocol that are not characterized in the second network protocol;

pruning the set of second network protocol FSMs to provide only operations corresponding to the services performed by the first complement service FSM to generate a first set of complement network protocol FSMs;

pruning the set of first network protocol FSMs to provide only operations corresponding to the

services performed by the second complement service FSM to generate a second set of complement network protocol FSMs; and

constructing a gateway system employing circuit equivalents of the first and second set of protocol FSMs, the first and second set of complement protocol FSMs and an interface converter, wherein each of the steps of pruning the set of protocol FSMs to generate a corresponding complement protocol FSM comprises the steps of:

deleting an edge from a first protocol service FSM corresponding to a service not characterized in a second protocol service FSM, and deleting an edge from the second protocol service FSM corresponding to a service not characterized in the first protocol service FSM to form a first set of pruned FSMs;

deleting an edge from the first set of pruned FSMs which represent service primitives for which there are no counterpart service primitives in the first set of pruned FSMs to form a second set of pruned FSMs; and

deleting an edge in the second set of pruned FSMs which are not within a strongly connected component of the corresponding second pruned FSM that contains the initial state to form a third set of pruned FSMs.

5. A method for generating a gateway system between first and second networks implementing corresponding first and second communication protocols which provide first and second sets of services, wherein each protocol can be represented by a set of at least one finite state machine (FSM), the method comprising the steps of:

determining a common subset of services equal to a subset of the intersection of services provided by the protocols;

determining a first complement service FSM for the first network protocol, wherein the first complement service FSM contains those services provided by the second network protocol that are not characterized in the first network protocol;

pruning the set of second network protocol FSMs to provide only operations corresponding to the services performed by the first complement service FSM to generate a first set of complement network protocol FSMs; and

constructing a gateway system employing circuit equivalents of the first and second set of protocol FSMs, the first set of complement protocol FSMs and an interface converter to provide all the input-output services of the first and second networks to users of the first network in communicating through the gateway system.

6. The method of claim 5 further comprising:
 - determining a second complement service

FSM for the second network protocol, wherein the second complement service FSM contains those services provided by the first network protocol that are not provided by the second network protocol;

pruning the set of first network protocol FSMs to provide only operations corresponding to the services performed by the second complement service FSM to generate a second set of complement network protocol FSMs, wherein the step of constructing the gateway system further includes employing a circuit equivalent of the second set of complement protocol FSMs and an interface converter to provide all the input-output services of the first and second networks to users of the first and second network in communicating through the gateway system.

7. A gateway system between first and second networks, each network implementing a protocol, the gateway system comprising:

a first end node connected to the first network, the end node having a protocol circuit and a first complement protocol circuit;

a first communication channel connected to the first end node;

a second end node connected to the second network, the second end node having a protocol circuit;

a second communications channel connected to the second network end node; and

a protocol converter connected to the first and second end nodes, the protocol converter having a first interface protocol circuit for communicating with the first end node protocol circuit, and a first interface complement protocol circuit for communicating with the first complement protocol circuit, the protocol converter further having a second interface protocol circuit for communicating with the second end node protocol circuit, and an interface converter connected to the first and second interface protocol circuits and first interface complement protocol circuit, wherein the interface converter translates and routes information between the first interface protocol circuit and the second interface protocol circuit, and between the first interface complement protocol circuit and the second interface protocol circuit to provide the input-output services of the first and second networks to users of the first network in communicating through the gateway system.

8. The gateway system of claim 7, wherein the circuits in an end node are contained in a single processor.

9. The gateway system of claim 7, wherein at least two of the circuits in the protocol converter are contained in a single processor.

10. The gateway system of claim 7, wherein the signals

between the first protocol circuit and the first interface protocol circuit and between the first complement protocol circuit and the first interface complement protocol circuit are multiplexed on the first communications channel.

11. The gateway system of claim 7 further comprising:
a second complement protocol circuit disposed in the second end node, the second complement protocol circuit being connected to the second network and the second communications channel; and

a second interface complement protocol circuit disposed in the protocol converter and being connected to the interface converter and the second communications channel, wherein the interface converter routes information between the first interface protocol circuit and the interface complement protocol circuit, and the second interface protocol circuit and second interface complement protocol circuit to provide input-output services of the first and second network protocols to users of the first and second networks in communicating through the gateway system.

12. The gateway system of claim 11, wherein the circuits in the second end node are contained in a single processor.

13. The gateway system of claim 11, wherein at two of the circuits in the protocol converter are contained in a single processor.

14. The gateway system of claim 11, wherein the signals between the second protocol circuit and second complement protocol circuit and the corresponding second interface protocol circuit and the second interface complement protocol circuit are multiplexed on the second communications channel, respectively.

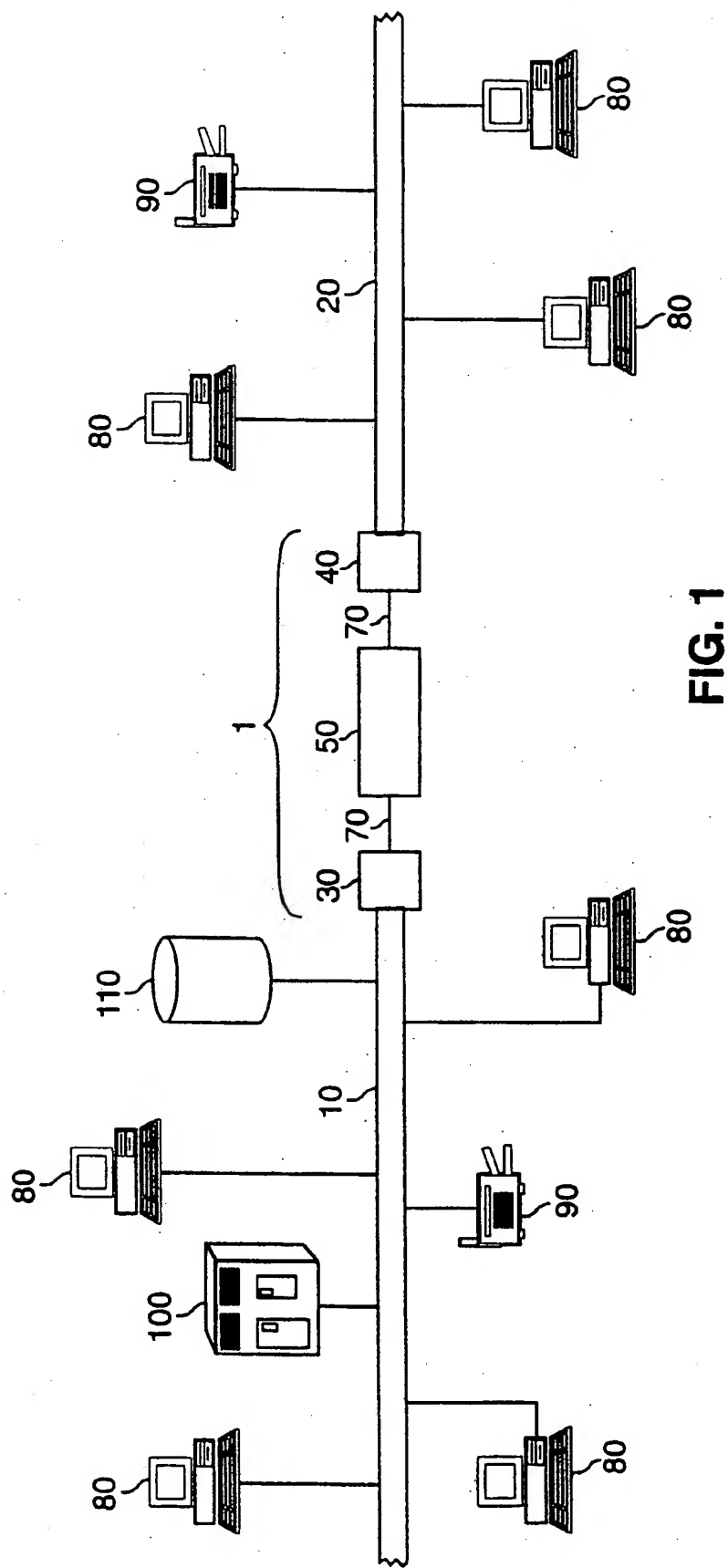


FIG. 1

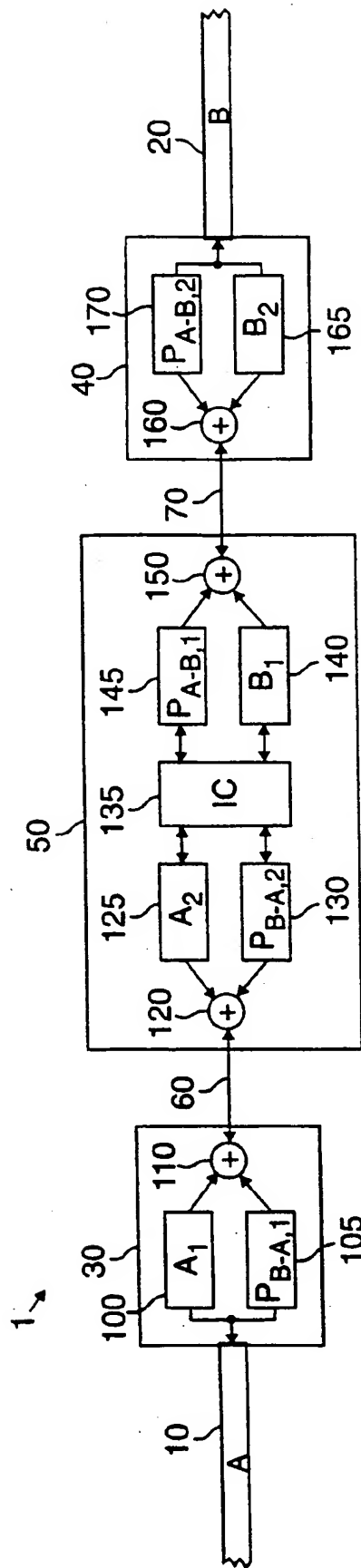


FIG. 2

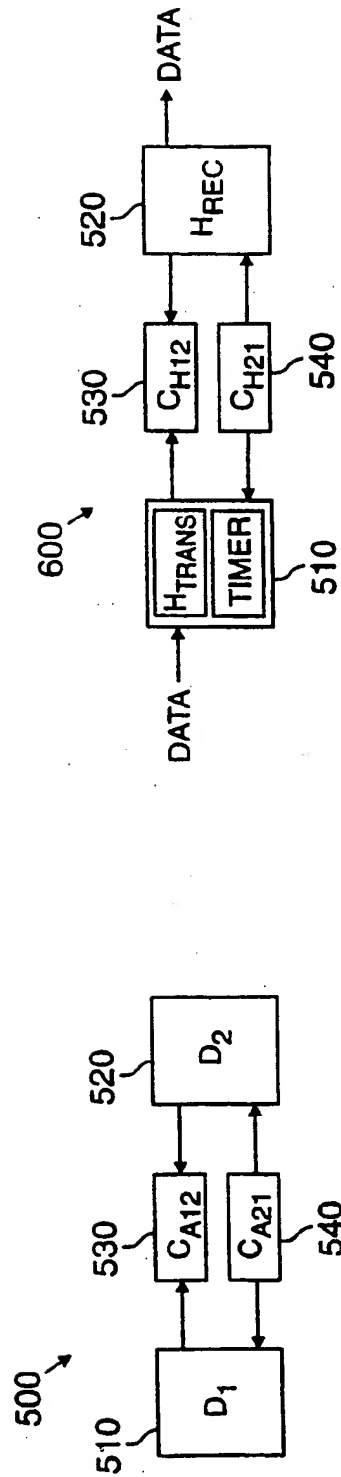
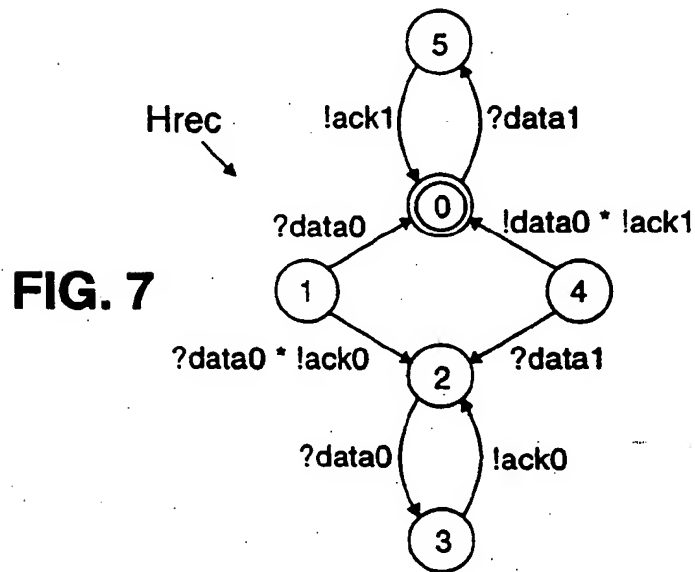
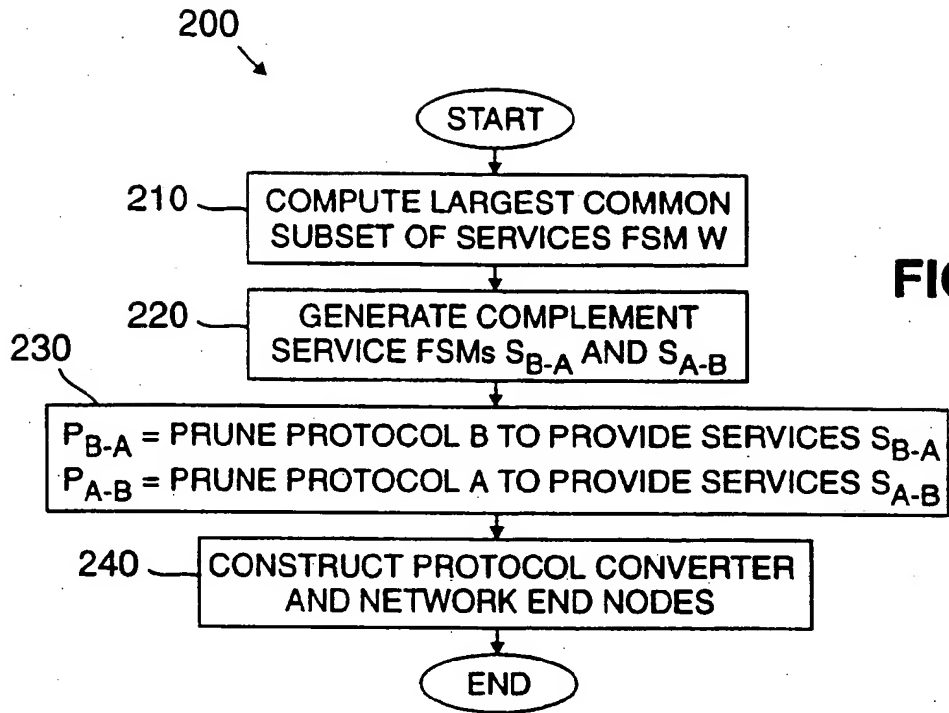


FIG. 5

FIG. 12



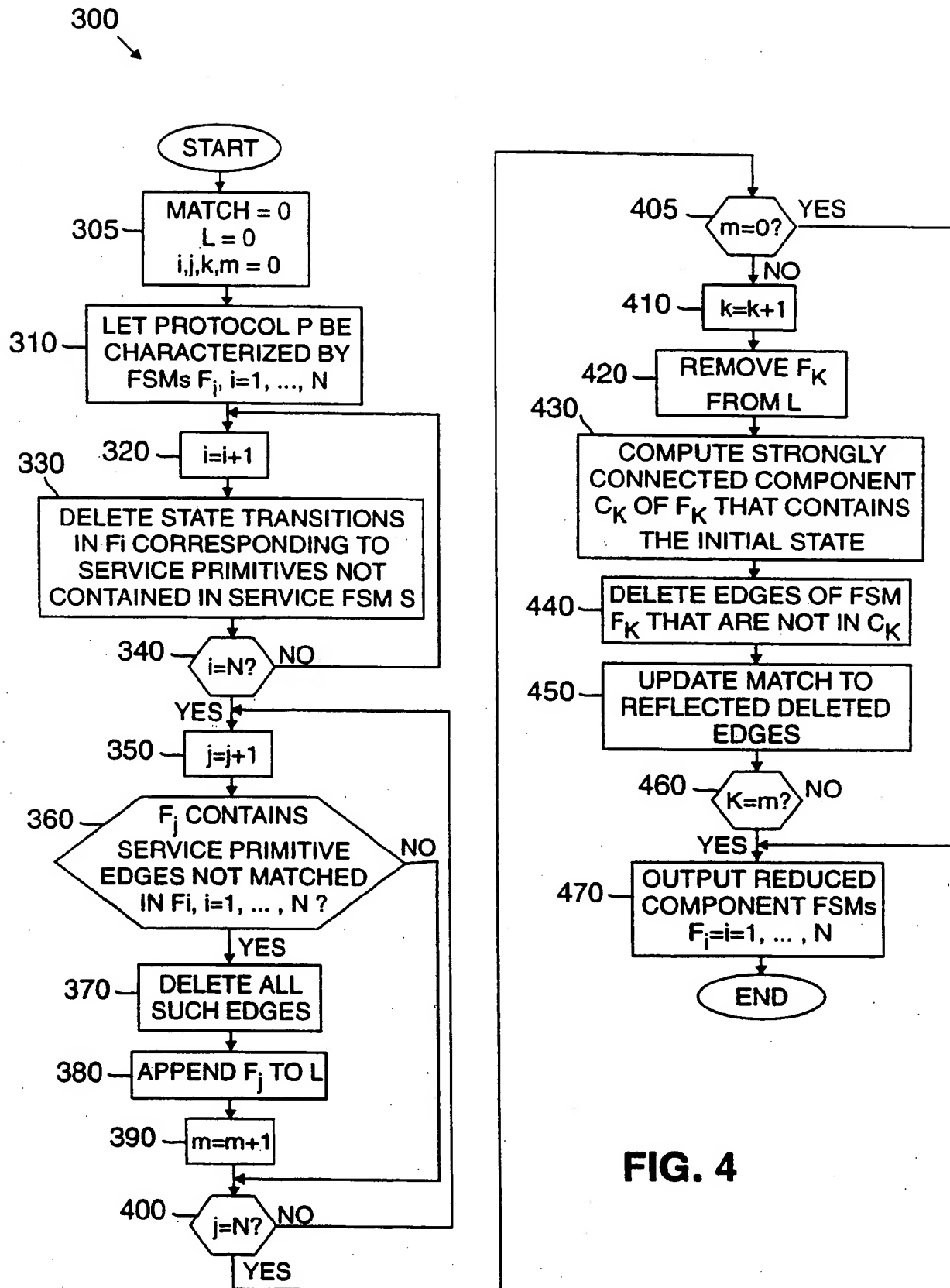


FIG. 4

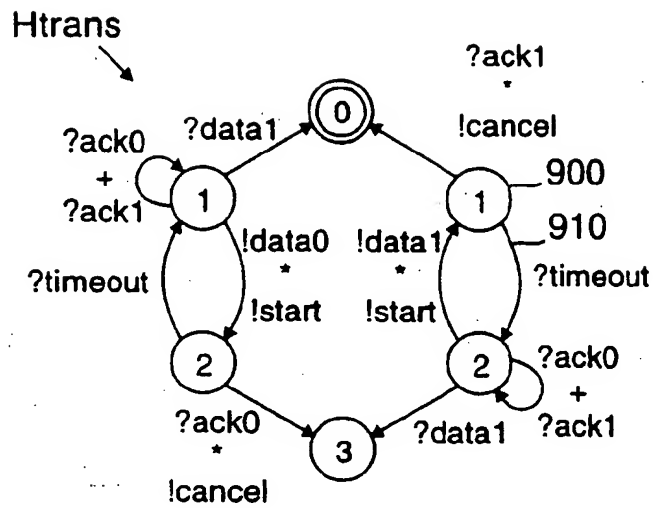


FIG. 8

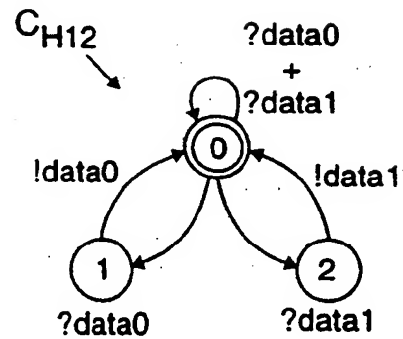
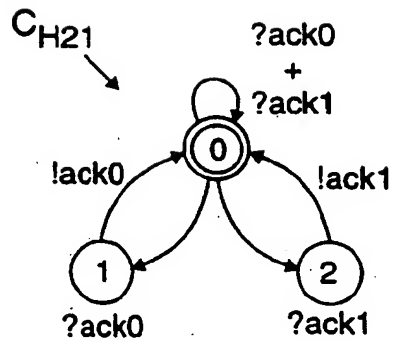


FIG. 9



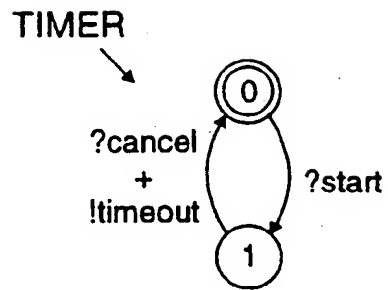


FIG. 10

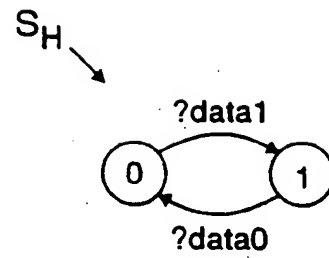


FIG. 11

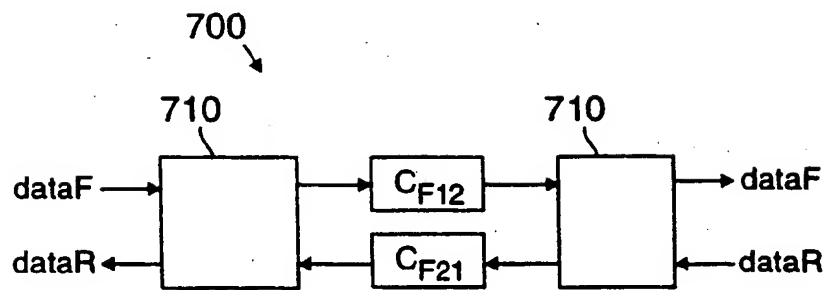


FIG. 13

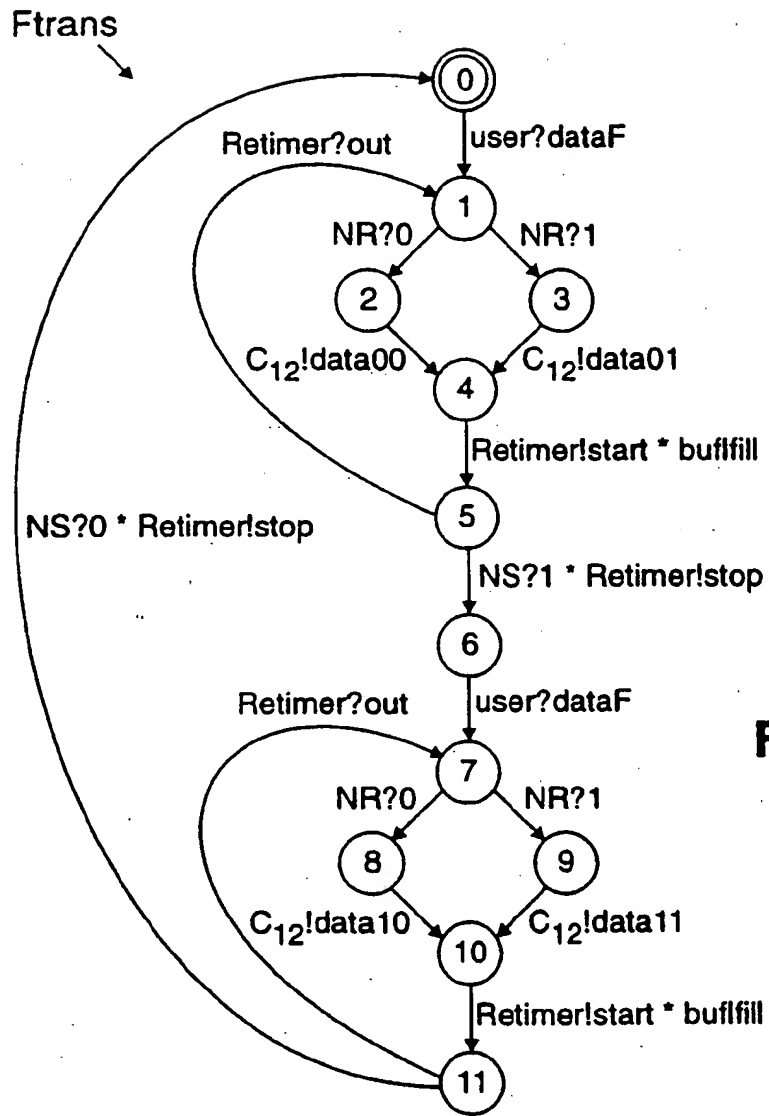


FIG. 14

FIG. 15

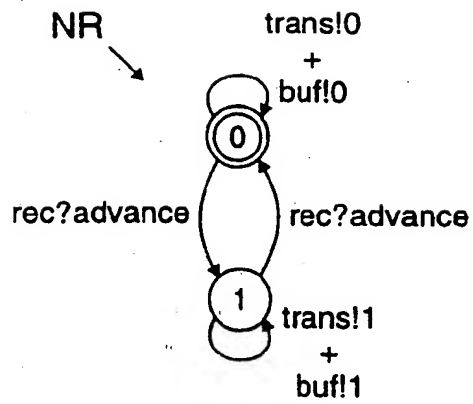
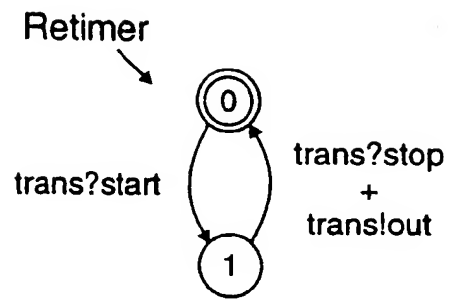
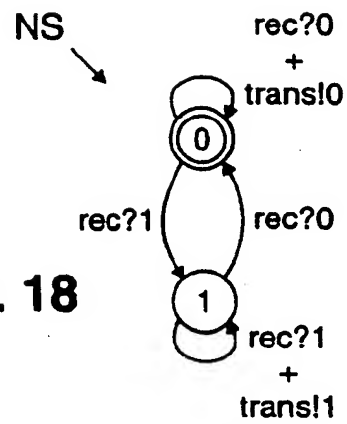


FIG. 17

FIG. 18



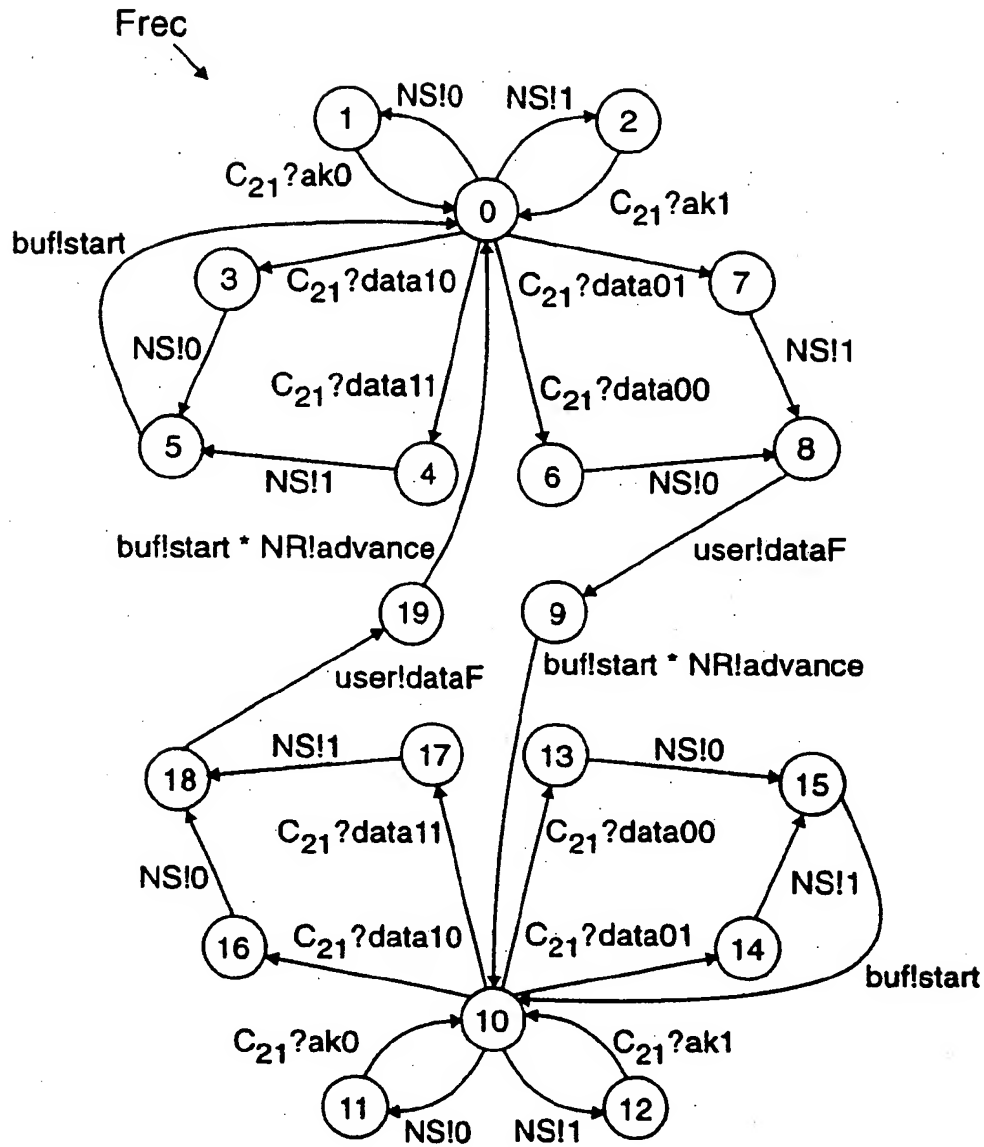
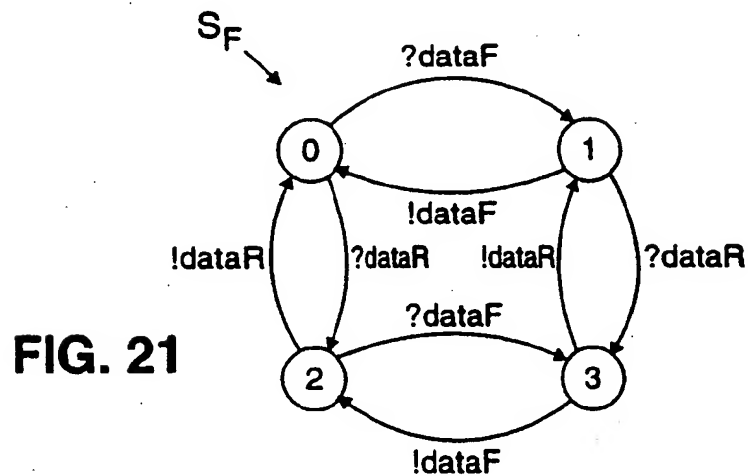
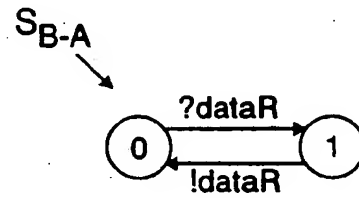
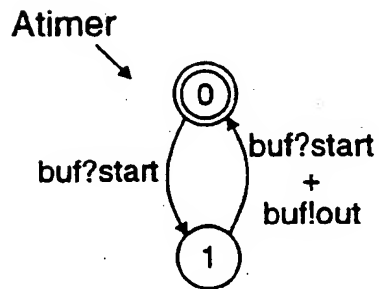
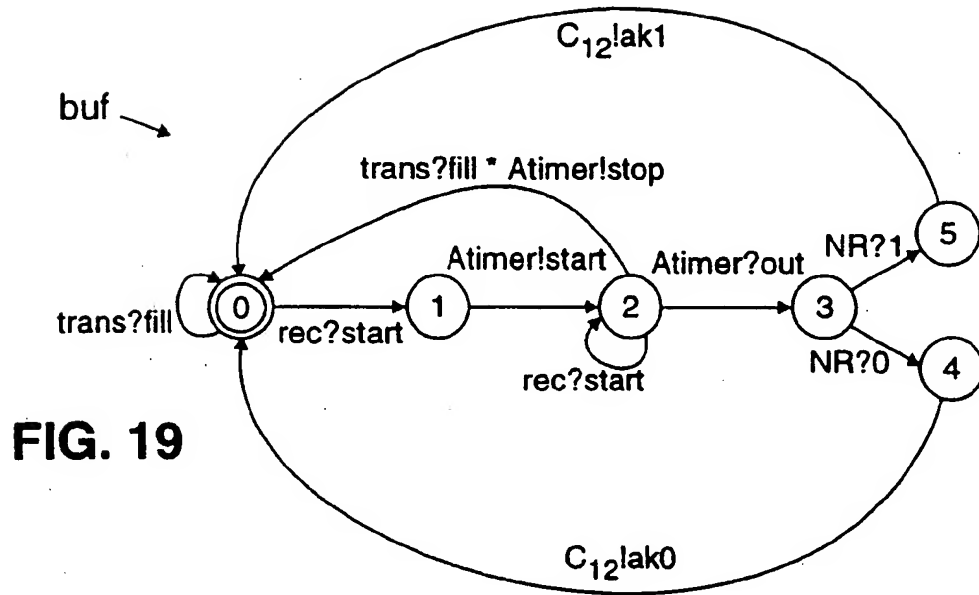


FIG. 16



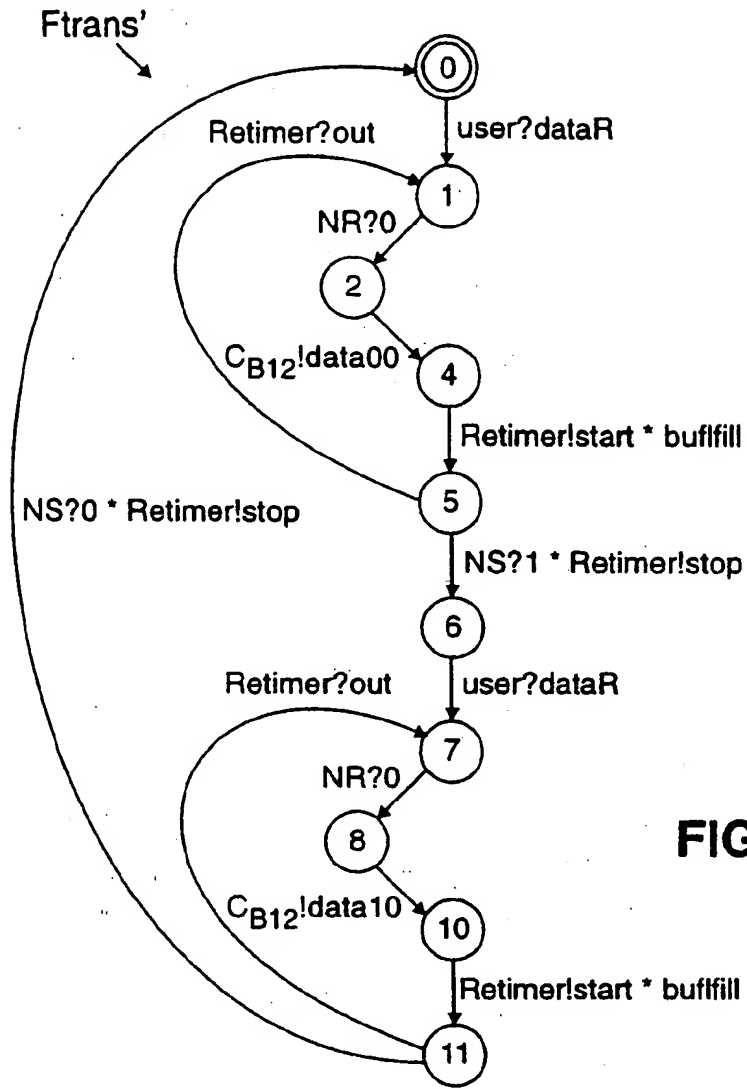


FIG. 23

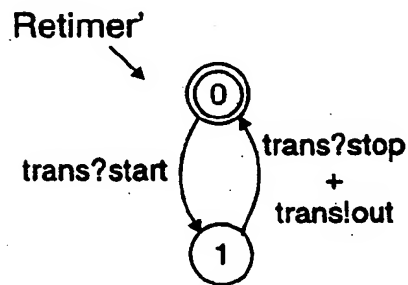


FIG. 24

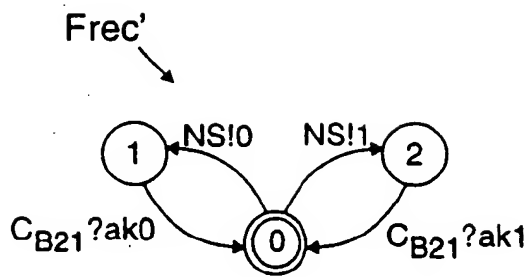


FIG. 25

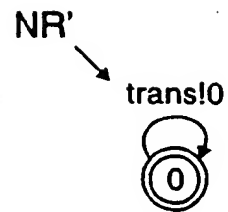


FIG. 26

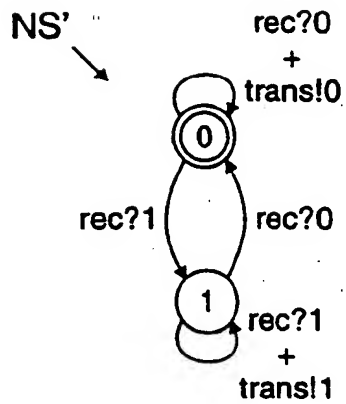


FIG. 27

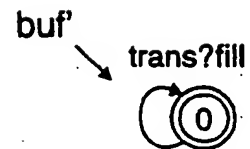


FIG. 28

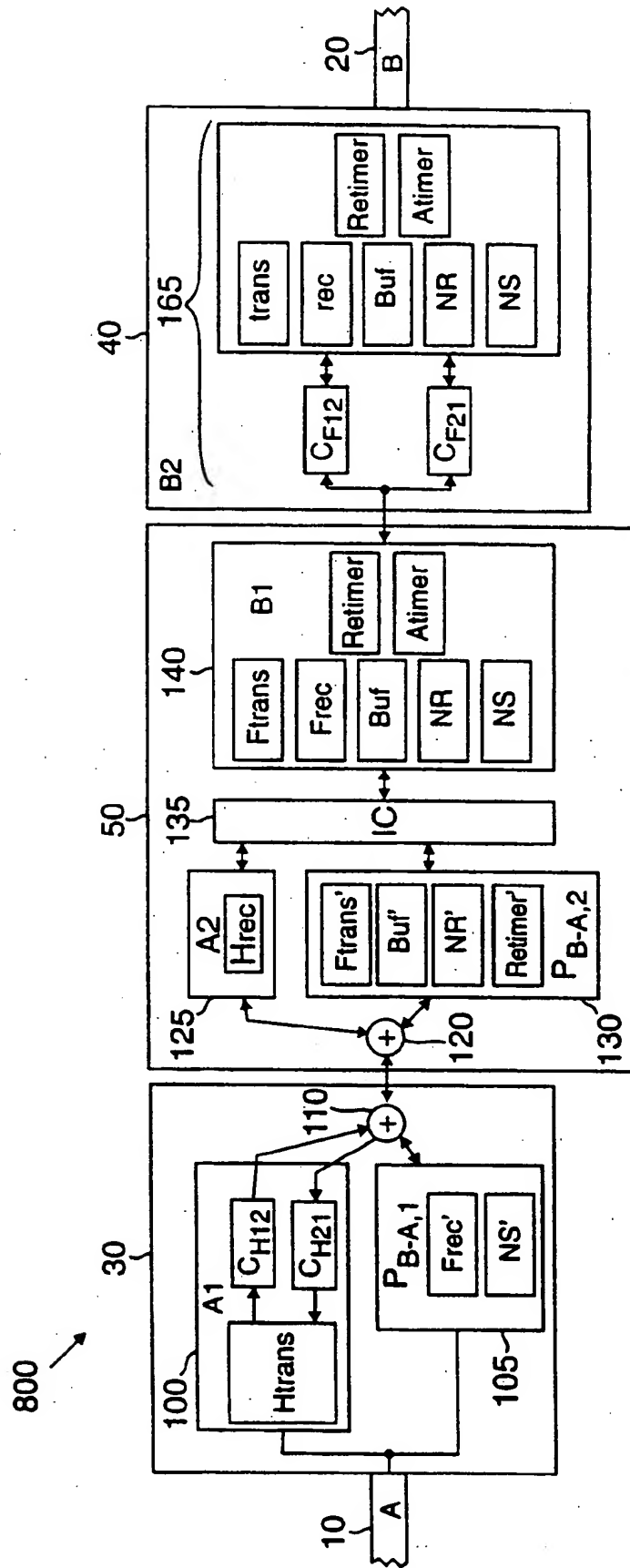


FIG. 29



EUROPEAN PATENT APPLICATION

(88) Date of publication A3:
29.03.2000 Bulletin 2000/13

(51) Int Cl.7: H04L 29/00, H04L 29/06

(43) Date of publication A2:
03.04.1996 Bulletin 1996/14

(21) Application number: 95306439.1

(22) Date of filing: 14.09.1995

(84) Designated Contracting States:
DE FR GB NL

• Sabnani, Krishan Kumar
Westfield, New Jersey 07090 (US)

(30) Priority: 20.09.1994 US 309213

(74) Representative:
Watts, Christopher Malcolm Kelway, Dr. et al
Lucent Technologies (UK) Ltd,
5 Mornington Road
Woodford Green Essex, IG8 0TU (GB)

(71) Applicant: AT&T Corp.
New York, NY 10013-2412 (US)

(72) Inventors:
• Netravali, Arun Narayan
Westfield, New Jersey 07090 (US)

(54) Computer network gateway

(57) A gateway system for connecting first and second communication networks has first and second network end nodes and a protocol converter. The gateway system operates to provide input-output services of both network protocols to users of both networks in communicating through the gateway system. The gateway system employs associated complement protocol circuits in the end nodes and the protocol converter to provide

these services. An efficient technique for generating the complement protocol of a first network is to prune finite state machines representing the operation of the second network protocol to provide the services of the second protocol that are not provided by the first protocol. The complement protocol circuits for the second network are then constructed based on the pruned finite state machines representing the operation of the first protocol.

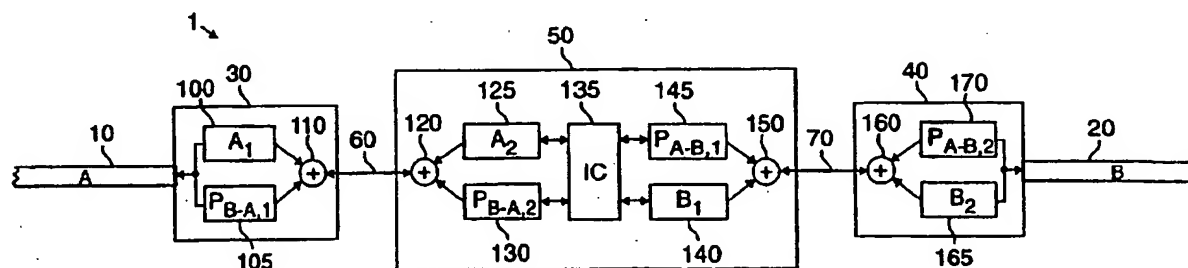


FIG. 2



European Patent
Office

EUROPEAN SEARCH REPORT

Application Number
EP 95 30 6439

DOCUMENTS CONSIDERED TO BE RELEVANT			
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int.Cl.8)
X	SAHA D ; DHAR P : "A fast protocol conversion technique using reduction of state transition graphs" ELEVENTH ANNUAL INTERNATIONAL PHOENIX CONFERENCE ON COMPUTERS AND COMMUNICATIONS (CAT. NO.92CH3129-4), 1 - 3 April 1992, pages 628-635, XP002129458 Scottsdale, AZ, USA	1-7,11	H04L29/00 H04L29/06
Y	* abstract *	8	
A	* figures 1,2,5,6 *	9,10	
	* paragraphs '0002!,'0004! *		
X	JANAKI AKELLA ET AL: "SYNTHESIZING CONVERTERS BETWEEN FINITE STATE PROTOCOLS" PROCEEDINGS OF THE INTERNATIONAL CONFERENCE ON COMPUTER DESIGN - VLSI IN COMPUTERS AND PROCESSORS,US,LOS ALAMITOS, IEEE. COMP. SOC. PRESS, vol. -, 1991, pages 410-413, XP000245923 ISBN: 0-8186-2270-9	1,2,4-6	
A	* the whole document *	3,7-11	TECHNICAL FIELDS SEARCHED (Int.Cl.6)
Y	KRISTOL D M: "A POLYNOMIAL ALGORITHM FOR GATEWAY GENERATION FROM FORMAL SPECIFICATIONS" IEEE / ACM TRANSACTIONS ON NETWORKING,US,IEEE INC. NEW YORK, vol. 1, no. 2, 1 April 1993 (1993-04-01), pages 217-229, XP000385471 ISSN: 1063-6692	8	H04L
A	* abstract *	1-7,9-11	
	* figures 1,3,6-10 *		
	* page 217, column 2, line 32 - page 218, column 1, line 28 *		
	* paragraphs '000A!-'000C! *		
	-/--		
The present search report has been drawn up for all claims			
Place of search THE HAGUE		Date of completion of the search 1 February 2000	Examiner Cichra, M
<p>CATEGORY OF CITED DOCUMENTS</p> <p>X : particularly relevant if taken alone Y : particularly relevant if combined with another document of the same category A : technological background O : non-written disclosure P : intermediate document</p> <p>T : theory or principle underlying the invention E : earlier patent document, but published on, or after the filing date D : document cited in the application L : document cited for other reasons & : member of the same patent family, corresponding document</p>			

EPO FORM 1503 03.92 (Pct/C01)



European Patent
Office

EUROPEAN SEARCH REPORT

Application Number

EP 95 30 6439

DOCUMENTS CONSIDERED TO BE RELEVANT			
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int.Cl.6)
A	EP 0 560 706 A (IBM) 15 September 1993 (1993-09-15) * abstract * * claim 1 * * figure 2 * * page 3, line 1 - page 5, line 24 *	1,4,5,7	
A	DHAR P ; GANGULY R C ; DAS S ; SAHA D : "Network interconnection and protocol conversion-a protocol complementation approach" 'TECHNOLOGY ENABLING TOMORROW' 1992 IEEE REGION 10 , 11 - 13 November 1992, pages 116-120, XP002129459 Melbourne * the whole document *	1,3,5,7	
			TECHNICAL FIELDS SEARCHED (Int.Cl.6)
The present search report has been drawn up for all claims			
Place of search THE HAGUE		Date of completion of the search 1 February 2000	Examiner Cichra, M
<p>CATEGORY OF CITED DOCUMENTS</p> <p>X : particularly relevant if taken alone Y : particularly relevant if combined with another document of the same category A : technological background O : non-written disclosure P : intermediate document</p> <p>T : theory or principle underlying the invention E : earlier patent document, but published on, or after the filing date D : document cited in the application L : document cited for other reasons & : member of the same patent family, corresponding document</p>			

EPO FORM 1503 03/82 (P04C01)

This annex lists the patent family members relating to the patent documents cited in the above-mentioned European search report. The members are as contained in the European Patent Office EDP file on The European Patent Office is in no way liable for these particulars which are merely given for the purpose of information.

FPO FOAM P0459

For more details about this annex : see Official Journal of the European Patent Office, No. 12/82